

Some Results on the Foundations of Elementary Plane Euclidean and Non-Euclidean Geometries

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Outline

1 Some History...

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- 2 Elementary Geometries
 - Hilbert planes
 - Pythagorean planes
 - Euclidean planes
 - Tarski's elementary Geometry

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- 3 Metamathematical Results
 - Arithmetic
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By “*elementary*” one means, in the words of Alfred Tarski:

“(...) that part of Euclidean geometry which can be formulated and established without the help of any set-theoretical devices.” [4]

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“Elementary Euclidean geometry is a much more ancient and simple subject than the axiomatic theory of real numbers (...)” [1]

As Robin Hartshorn [3] said, [without real numbers] *“the true essence of geometry can develop most naturally and economically.”*

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Five groups of axioms:

- I Incidence
- II Betweenness (or Order)
- III Congruence
- IV Parallelism axiom
- V Continuity axiom

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The study of Hilbert planes allows for the study of straightedge and compass constructions in plane geometry (Greenberg [1]).

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For Hilbert planes, this axiom is equivalent to Euclid's fifth postulate.

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If \mathcal{P} is a Pythagorean plane then it's possible, by fixing a line ℓ and a point “ O ” on that line, to define an ordered field F .

That field will be called a Pythagorean field since for all $a, b \in F$
 $\sqrt{a^2 + b^2} \in F$.

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If we now consider the Cartesian plane F^2 , we will have $F^2 \simeq \mathcal{P}$.

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LINE-CIRCLE AXIOM

If a line passes through a point inside a circle, then it intersects the circle in two distinct point.

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Euclid I.1

Euclid I.1 does not hold in all Hilbert planes.

(It does, however, hold in all Pythagorean planes, by a different construction, using the fact that the Pythagorean field $F \ni \sqrt{3}$.)

In Hilbert planes, Circle-Circle, Line-Circle and the Triangle Theorem are all equivalent.

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Any model of I, II, III, IV and the Circle-Circle Axiom is called an **Euclidean plane**.

Any Euclidean plane is isomorphic to a Cartesian plane F^2 , where F is some arbitrary *Euclidean field* – an ordered field in which every positive element has a square root, i.e.,
 $\forall a \in F (a > 0 \rightarrow \exists b \in F b^2 = a)$.

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Every plane geometric proposition in Euclid's *Elements* can be proved from I, II, III, IV and the Circle-Circle Axiom.

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He defined the relation of “collinearity” of three points in terms of β , and so he did not need “line” as a primitive notion:

$$\text{“}x, y \text{ and } z \text{ are collinear”} \Leftrightarrow \beta(x, y, z) \vee \beta(x, z, y) \vee \beta(y, x, z)$$

If we consider

DEDEKIND'S CUT AXIOM (V)

$$\forall X, Y (\exists z \forall x, y (x \in X \wedge y \in Y \rightarrow \beta(z, x, y)) \\ \rightarrow \exists u \forall x, y (x \in X \wedge y \in Y \rightarrow \beta(x, u, y)))$$

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ELEMENTARY AXIOM SCHEMA OF CONTINUITY $((V)_{\text{elem}})$

$$\exists z \forall x, y (\varphi(x) \wedge \psi(y)) \rightarrow \beta(z, x, y) \\ \rightarrow \exists u \forall x, y (\varphi(x) \wedge \psi(y)) \rightarrow \beta(x, u, y)$$

for every formulas φ and ψ

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Hilbert planes \supset Pythagorean planes \supset Euclidean planes \supset
 \supset Tarski-elementary Euclidean planes \supset the real Euclidean plane

Arithmetic

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Recall that the Second-Order Arithmetic theory (in the language with $0, S, +, \cdot, =, <$ and \in), contains as one of its axioms

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$$\forall X \left(0 \in X \wedge \forall x (x \in X \rightarrow S(x) \in X) \right) \rightarrow \forall x (x \in X)$$

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and that if we replace it by its “elementary” axiom schema of induction,

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we end up with an axiom system for First-Order Peano Arithmetic.

Definition

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It is well known that PA is **incomplete** – there exists a sentence φ such that neither φ nor $\neg\varphi$ is provable in PA – and **essentially undecidable** – i.e., any consistent extension of PA is undecidable. In particular, the set $\{\varphi \mid \varphi \text{ is true in } \mathbb{N}\}$ is undecidable.

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Also...

(K.Gödel) No finitary proof of the consistency of PA is possible.

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[Recall that every Tarski-elementary Euclidean plane is isomorphic to F^2 , where F is some real-closed ordered field.]

As Descartes showed, every geometric statement φ about the plane F^2 translates into an algebraic statement φ^* about F , and so:

$$TEG \vdash \varphi \leftrightarrow RCOF \vdash \varphi^*,$$

where TEG denotes the Tarski's Elementary Geometry theory and RCOF the theory of real-closed ordered fields.

Theorem (Tarski)

*RCOF is a **complete** theory.*

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Fact 1. RCOF is a **decidable** theory.

Fact 2. $\{\varphi \in \mathcal{L}_{RCOF} \mid \mathbb{R} \models \varphi\}$ is decidable.

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The theory of the Tarski-elementary Euclidean planes is **complete**, **decidable** and has a finitary proof of consistence.

Final Remarks

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Ziegler proved that any finitely axiomatizable first-order theory of fields, having the real number field \mathbb{R} as a model must be undecidable.

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Theorem

*The theory of Euclidean fields is **undecidable**.*

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$$EPG \vdash \varphi \leftrightarrow EF \vdash \varphi^*,$$

where EPG denotes the Euclidean plane Geometry theory and EF the theory of Euclidean fields.

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It follows that the Euclidean plane Geometry is **undecidable**:

“Elementary Euclidean geometry is genuinely creative, not mechanical” [1]

References

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