

THE MATHEMATICS OF MAKING A MESS

(AN INTRODUCTION TO RANDOM WALK ON GROUPS)

(NOTES FOR A TALK IN LISBON, OCT 13, 2020)

1) THE BASIC SET UP

LET G BE A FINITE GROUP

LET $S = S^{-1}$ BE A GENERATING SET (i.e. $\langle S \rangle = G$)

A RANDOM WALK ON G , STARTS AT e AND

$$X_{n+1} = \epsilon_{n+1} X_n$$

WITH ϵ_i i.i.d FROM

$$P(g) = \begin{cases} 1/|S| & \text{if } g \in S \\ 0 & \text{—} \end{cases}$$

$$P * P(g) = \sum_{h \in G} P(h) P(h^{-1}g), \quad P^k = P * P^{k-1}$$

$$U(g) = \frac{1}{|G|}$$

2) THEOREM (POINCARÉ = 1890)

$$P^k(g) \xrightarrow[\omega]{} U(g)$$

MORE GENERALLY, IF $\text{supp } P$ IS NOT IN A COSET OF A SUBGROUP

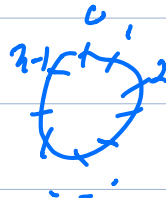
3) HOW FAST?

$$\|P^h - u\| = \max_{A \subseteq G} |P^h(A) - u(A)| = \frac{1}{2} \sum_{g \in G} |P^h(g) - u(g)|$$

GIVEN $\epsilon > 0$, HOW LARGE h SO $\|P^h - u\| < \epsilon$?

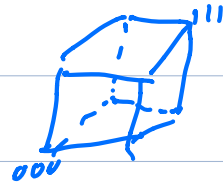
4) EXAMPLES

(a) $G = \mathbb{Z}_n$, $S = \{0, 1, -1\}$



ORDER n^2 STEPS ARE $n + S$.

(b) $G = \mathbb{Z}_2^n$, $S = \{0, e_1, \dots, e_n\}$



ORDER $d \log d$ STEPS $n + S$

(EVEN BETTER URM)

(c) $G = S_n$ COIN PERMUTATIONS

$S = \{\text{ALL TRANSPOSITIONS}\}$

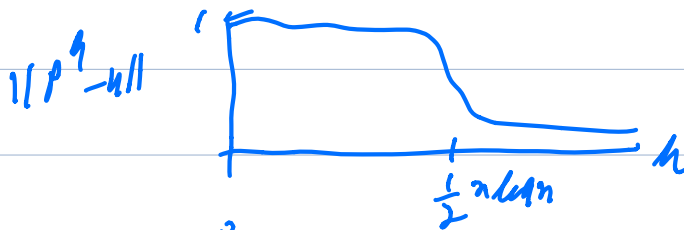
5) THEOREM (WITH SAHSHAHANI) ON S_n , LET

$$P(\sigma) = \begin{cases} 1/n & \sigma = id \\ 2/n^2 & \sigma = (i, j) \\ \vdots & \vdots \end{cases}$$

LET $h = \frac{1}{2} n \log n + c n$

$\|P^h - u\| \leq 2 e^{-2c}$ (MATCHING LOWER BOUND)

6) THERE IS A CUTOFF AT $\frac{1}{2} n \log n$



DEFINITION $\{G_n, P_n\}_{n=1}^{\infty}$ GROUPS

P_n HAS A CUTOFF AT k_n IF, $\forall \epsilon > 0$

$$\|P_n^{k_n(1+\epsilon)} - U\| \rightarrow 0$$

$n \uparrow \infty$

$$\|P_n^{k_n(1-\epsilon)} - U\| \rightarrow 1$$

CONJECTURE \forall SEQUENCE OF GENERATING SETS OF S_n THERE IS CUTOFF.

(INDEED \forall SEQUENCE OF FINITE SIMPLE GROUPS)

2) FOURIER ANALYSIS

A REPRESENTATION OF G IS $\rho: G \rightarrow GL(V)$

$$\rho(AB) = \rho(A)\rho(B)$$

eg

$$G = \mathbb{Z}_n \quad \rho_j(g) = e^{2\pi i j h / n}$$

$$G = \mathbb{Z}_2^n \quad \rho_x(y) = (-1)^{x \cdot y}$$

$$G = S_n, \quad \rho_{\text{triv}}(\sigma) = 1, \quad \rho_{\text{sgn}}(\sigma) = \text{sgn}(\sigma), \quad \rho_{\pi}(\sigma) \in \{-1, 1\} \text{ PERMUTATION MATRICES}$$

$$d_{\rho} = \dim V \text{ - DIMENSION OF } \rho$$

ρ IS IRREDUCIBLE IFF $\exists W \subseteq V$, NON TRIVIAL

$\rho|_W \subseteq W$ ALL $g \in G$.

LET P BE A PROBABILITY ON G .

$$\hat{P}(\rho) = \sum_g P(g) \rho(g)$$

FACTS . $P * P(\rho) = \hat{P}(\rho)^2$

. $\hat{u}(\rho) = \begin{cases} 0 & \rho \text{ IRREDUCIBLE, NOT TRIVIAL} \\ 1 & \rho = \text{TRIV.} \end{cases}$

. $P^k(g) = \frac{1}{|G|} \sum_{\rho \in \hat{G}} d_\rho \text{Tr}(\rho(g^{-1}) \hat{P}(\rho)^k)$

SO SHOW $P^k \Rightarrow u$ VIA $\hat{P}(\rho)^k \rightarrow 0$.

LEMMA (UPPER BOUND LEMMA) $4 \|P^k - u\|^2 \leq \sum_{\rho \neq 1} d_\rho \|\hat{P}(\rho)\|^k$ (TRACE NORM)

PROOF $4 \|P^k - u\|^2 = (\sum_g |P(g) - u(g)|)^2 \leq |G| \sum_g |P(g) - u(g)|^2$
 $= \sum_{\rho \neq 1} d_\rho \|\hat{P}(\rho)\|^k$ (PLANCHEREL).

8) BACK TO RANDOM TRANSPOSITIONS ON S_n

KEY FACT P IS CONSTANT ON CONJUGACY CLASSES

$$P(\bar{\sigma} \sigma s) = P(\sigma)$$

SCHUR'S LEMMA (ANY SUCH G ON ANY SUCH GROUP)

$$\hat{P}(\rho) = c I, \quad c = \text{Tr} \hat{P}(\rho) / d_\rho$$

eg. $P(\sigma) = \begin{cases} 1/n & \sigma = 1 \\ 2/n^2 & \sigma = (i,j) \\ 0 & - \end{cases}$

$$\hat{P}(k) = c I \quad c = \left(\frac{1}{n} + \frac{n-1}{n} \frac{\chi_e(12)}{d_e} \right)$$

$$\chi_e(\sigma) = \text{Tr}(P(\sigma)) \text{ CHARACTER}$$

so

$$\|P^k - U\|^2 \leq \sum_{\substack{e \in S_n \\ e \neq 1}} d_e \left(\frac{1}{n} + \frac{n-1}{n} \frac{\chi_e(12)}{d_e} \right)^{2k}$$

9) FINALLY, WE'VE GONE AS FAR AS WE CAN "WITHOUT KNOWING ANYTHING"

- WHAT IS \hat{G}
- WHAT IS d_e
- WHAT IS $\chi_e(12)$

LOOK AT CHARACTER TABLES, USUALLY $\chi_e(12)/d_e$ SMALL ($= 1/2$)

IF SO

$$\sum_e d_e \left(\frac{1}{n} + \frac{n-1}{n} \frac{\chi_e}{d_e} \right)^{2k} \approx \sum_e d_e \left(\frac{1}{2} \right)^{2k} = \frac{n!}{2^{2k}}$$

so $k \gg n \log n$ MAKES THIS SMALL.

BUT $n-1$ DIMENSIONAL REPRESENTATION

$$d_{n-1} = n-1, \quad \chi_e = n-3, \quad \left(\frac{1}{n} + \frac{n-1}{n} \frac{\chi_e}{d_e} \right)^{2k} = \left(1 - \frac{2}{n} \right)^{2k}$$

HAVE TO MAKE

$$(n-1)^2 \left(1 - \frac{2}{n} \right)^{2k} \leq e^{2k \log n - \frac{4k}{n}} \text{ Small}$$

$$k = \frac{1}{2} n \log n + cn \text{ MAKES IT } \leq e^{-4c}$$

10) COMPARISON THEORY

SAY YOU HAVE A GROUP G , GENERATING SET S

ALSO HAVE A GENERATING SET \tilde{S} AND

'KNOW EVERYTHING' ABOUT \tilde{S}

THEN, IF UNDERSTAND HOW TO WRITE ELTS OF S WITH \tilde{S}

YOU CAN GET GOOD ESTIMATES FOR P .

EXAMPLE $G = S_n$, $S = \{(12) (12 \dots n)\}$ VERY DIFFERENT FROM $\{(i,j)\}$.

$$\text{GET } \frac{1}{6} e^{-k/n^3 \log n} \leq \|P^n - U\| \leq 6 e^{-k/n^3 \log n}$$

INDEED THIS RESULT HOLDS FOR ALMOST EVERY PAIR $\{a, b\} \subseteq S_n$!

(HELFGOT, SERNOBY, ZUK) (WE CONSECUTIVE ZUKOFF)

11) PROGRAM

- FIND A (FAMILY OF) GROUPS YOU LIKE
- PICK A GENERATING SET OF INTEREST S
- FIND A SMALL GENERATING CONJUGACY CLASS (OR SET OF THEM) \tilde{S}
- DO THE WORK REQUIRED TO UNDERSTAND P_S
- DO THE WORK TO UNDERSTAND THE 'GEOMETRY OF THE Cayley GRAPH OF (G, S) IN TERMS OF \tilde{S} '
- SEND ME A COPY!

12) PROGRESS PART 1 IS IN PRETTY GOOD SHAPE (CONJUGACY)
(WORK OF GLUCK, LIEBECK, SHREVE, GURALNICK, TIEB, ...)
FOR FINITE GROUPS OF LIE TYPE.

$GL_n(q), U_n(q), \dots$ THERE ARE EVEN SOME
CUTOFF RESULTS (NAK + N+S)

BUT PART 2 REMAINS TO BE DONE
SPECIFIC PROBLEM: ALL SUCH GROUPS ARE
GENERATED BY TWO GENERATORS. THE
STEINBERG'S GENERATORS + GO FORWARD

(b) FOR LOW-CLASS NILPOTENT GROUPS (BAD # GENERATORS, NAK+S)
eg HERSENBURG GROUP $\begin{pmatrix} 1 & 4 & 3 \\ 0 & 1 & 4 \\ 0 & 0 & 1 \end{pmatrix}$ $4, 3 \in \mathbb{Z}_m$
(DIMENSION)² N+S (SALUF-COSTE)

FOURIER ANALYSIS

- POSSIBLE (CARLOS AMARÉ - SUPERCOMPUTER)
- DIFFICULT

(c) P-GROUPS OF MAXIMAL CLASS, HIGH COCLASS
JONATHAN HEARMAN, SAM THOMAS

(d) SOLVABLE GROUPS, eg $x \mapsto ax+b \pmod{p}$

- FASCINATING (HEDGECOCK, US, BAQUINAB, VARSY, EBERHARDT...)
- LARGELY OPEN.

(e) THERE IS PROGRESS: PROVING CUTOFF ON S_n FOR

- TOP TO RANDOM (WITH BLOCKS)
- RANDOM TO RANDOM (BERNSTEIN, WESTERLIND)

- ADJACENT TRANSPOSITIONS (LACROIX)
- BERNOULLI-LINDLIE,
- GOLDFORD PRIMS

(3) OTHER TECHNIQUES (DODOLUB, STATIONARY TIMES, LUG-SOBOLEV, GARCH THEORY (BEASTIKI, SCHARM), METHOD OF MIRACLES

(4) WE WORK ON ALL KINDS OF GROUPS; CONTINUOUS, DISCRETE, LIE, P-ADIC, ON INFINITE GROUPS RESULTS 'RATES OF CONVERGENCE' WITH 'RATE OF DECA' OF P^1 (LID), DIFFERENT METRICS

(5) IN ADDITION TO 'RATES OF CONVERGENCE' WE DO

- FIRST HITTING TIME (TO POINT, SET)
- COVERTIMES (SIZE OF RANGE)
- BIRTHDAY PROBLEMS
- BEHAVIOUR BEFORE CONVERGENCE

(6) REAL PROBLEMS: ALL OF ABOVE IS 'WARM UP' FOR AN ACTUALLY IMPORTANT PROBLEM;

TAKE ANY ACTUAL MARKOV CHAIN USED BY A PHYSICIST, CHEMIST, BIOLOGIST, STATISTICIAN AND TRY TO PROVE ANYTHING USEFUL ABOUT ITS MIXING TIME.

17)

REFERENCES

LEVIN, D. AND PERES, Y. 'MARKOV CHAINS AND MIXING TIMES'
(AMS & FASE ON DAVID LEVIN'S HOME PAGE)

DIACONIS, P. 'THE MARKOV CHAIN MONTE CARLO REVOLUTION'
(Bull. AMS & ON MY HOME PAGE)

RANDOM WALKS ON GRAPHS - CHARACTERS AND GEOMETRY.

(EASIEST TO FIND ON MY HOME PAGE)

SALOFF-ZIGSB, L. PROBABILITY ON GRAPHS: RANDOM
WALKS AND INVARIANT DIFFUSIONS (NOTICES AMS)

+ I'M STILL HARD AT WORK ON THIS
PLEASE LOOK AT MY HOME PAGE
FOR LATEST !!!